# Cs502 Solved MCQS for Final Term <br> Without Repetitions 

## Question \# 1

Word Algorithm comes from the name of the Muslim author:
Abu Ja'far Mohammad ibn Musaal-Khowarizmi. (p7)

## Question \# 2

Al-Khwarizmi's work was written in a book titled
al Kitab al-mukhatasar fi hisab al-jabr wa'l-muqabalah (p7)
Question \# 3
The running time of quick sort depends heavily on the selection of:
No of inputs
Arrangement of elements in array
Size of elements
Pivot element (p49)
Question \# 4
Which sorting algorithm is faster?
$\mathrm{O}\left(\mathrm{n}^{\wedge} 2\right)$
O(nlogn)
$\mathrm{O}(\mathrm{n}+\mathrm{k})$
$\mathrm{O}\left(\mathrm{n}^{\wedge}\right)$

## Question \# 5

There is relationship between number of back edges and number of cycles in DFS Select correct option:
Both are equal.
Cycles are half of back edges.
Cycles are one fourth of back edges.
There is no relationship between back edges and number of cycle (p131)

## Question \# 6

You have an adjacency list for G , what is the time complexity to compute Graph transpose $\mathrm{G}^{\wedge} \mathrm{T}$.?
Select correct option:


## Question \# 7

Dijkstra's algorithm:
Has greedy approach to find all shortest paths
Has both greedy and dynamic approach to find all shortest paths
Has greedy approach to compute single source shortest paths to all other vertices (p154)
Has both greedy and dynamic approach to compute single source shortest paths to all other vertices.
Question \# 8

What is the time complexity to extract a vertex from the priority queue in Prim's algorithm?
Select correct option:
O (log E)
(V)
(V+E)
O $(\log V)$

## Question \# 9

Which is true statement in the following?
Kruskal algorithm is multiple source technique for finding MST.
Kruskal's algorithm is used to find minimum spanning tree of a graph, time complexity of this algorithm is $\mathrm{O}(\mathrm{EV})$
Both of above
Kruskal's algorithm (choose best non-cycle edge) is better than Prim's (choose best tree edge) when the graph has relatively few edges.

## Question \# 11

Kruskal's algorithm (choose best non-cycle edge) is better than Prim's (choose best tree edge) when the graph has relatively few edges.
True
False

## Question \# 12

Q knapsack problem is called a " $0-1$ " problem, because
Each item must be entirely accepted or rejected
(p92)

## Question \# 13

Which statement is true?
Select correct option:
If a dynamic-programming problem satisfies the optimal-substructure property, then a locally optimal solution is globally optimal.
If a greedy choice property satisfies the optimal-substructure property, then a locally optimal solution is globally optimal.
Both of above (p77, p98)
None of above
Question \# 14
A dense undirected graph is:
Select correct option:
A graph in which $\mathrm{E}=\mathrm{O}\left(\mathrm{V}^{2}\right) \quad$ (Reference)
A graph in which $\mathrm{E}=\mathrm{O}(\mathrm{V})$
A graph in which $\mathrm{E}=\mathrm{O}(\log \mathrm{V})$
All items above may be used to characterize a dense undirected graph

## Question \# 15

What algorithm technique is used in the implementation of Kruskal solution for the MST?
Greedy Technique
(p142)
Divide-and-Conquer Technique
Dynamic Programming Technique
The algorithm combines more than one of the above techniques

## Question \# 16

A digraph is strongly connected under what condition?
A digraph is strongly connected if for every pair of vertices $u, v e V, u$ can reach $v$.
A digraph is strongly connected if for every pair of vertices $u, v e V$, $u$ can reach $v$ and vice versa. (p135)
A digraph is strongly connected if for at least one pair of vertex $u$, $v e V, u$ can reach $v$ and vice versa.
A digraph is strongly connected if at least one third pair of vertices $u$, $v e V, u$ can reach $v$ and vice versa.

## Question \# 17

The relationship between number of back edges and number of cycles in DFS is,
Both are equal
Back edges are half of cycles
Back edges are one quarter of cycles
There is no relationship between no. of edges and cycles

## Question \# 18

Suppose that a graph $\mathrm{G}=(\mathrm{V}, \mathrm{E})$ is implemented using adjacency lists. What is the complexity of a breadth-first traversal of G? Select correct option:
$\mathrm{O}\left(|\mathrm{V}|^{\wedge} 2\right)$
$\mathrm{O}(|\mathrm{V}||\mathrm{E}|)$
$\mathrm{O}(|\mathrm{V}| \wedge 2|\mathrm{E}|)$
$\mathbf{O}(|\mathbf{V}|+|\mathbf{E}|)$

## Question \# 19

Forward edge is?
Select correct option:
$(u, v)$ where $u$ is a proper descendent of $v$ in the tree.
$(u, v)$ where $v$ is a proper descendent of $u$ in the tree.
(p129)
$(u, v)$ where $v$ is a proper ancestor of $u$ in the tree.
$(u, v)$ where $u$ is a proper ancestor of $v$ in the tree.

## Question \# 20

Back edge is?
Select correct option:
$(u, v)$ where $v$ is an ancestor of $u$ in the tree.
$(u, v)$ where $u$ is an ancestor of $v$ in the tree.
$(u, v)$ where $v$ is an predecessor of $u$ in the tree.
None of above

## Question \# 21

Cross edge is?
$(u, v)$ where $u$ and $v$ are not ancestor of one another
$(u, v)$ where $u$ is ancestor of $v$ and $v$ is not descendent of $u$.
$(u, v)$ where $u$ and $v$ are not ancestor or descendent of one another ( $\mathbf{p 1 2 9 )}$
$(u, v)$ where $u$ and $v$ are either ancestor or descendent of one another.

## Question \# 22

If you find yourself in maze the better traversel approach will be?
Select correct option:

## BFS

BFS and DFS both are valid

Level order
DFS
Question \# 23
In digraph $\mathrm{G}=(\mathrm{V}, \mathrm{E})$; G has cycle if and only if
Select correct option:
The DFS forest has forward edge.
The DFS forest has back edge (p131)
The DFS forest has both back and forward edge
BFS forest has forward edge

## Question \# 24

What is generally true of Adjacency List and Adjacency Matrix representations of graphs?
Select correct option:
Lists require less space than matrices but take longer to find the weight of an edge (v1,v2)
Lists require less space than matrices and they are faster to find the weight of an edge ( $\mathrm{v} 1, \mathrm{v} 2$ ) ( p 116 )
Lists require more space than matrices and they take longer to find the weight of an edge ( $\mathrm{v} 1, \mathrm{v} 2$ )
Lists require more space than matrices but are faster to find the weight of an edge ( $\mathrm{v} 1, \mathrm{v} 2$ )
http://cs-mcqs.blogspot.com/2012/06/data-structures-algorithms-multiple.html

## Question No: 25

Although it requires more complicated data structures, Prim's algorithm for a minimum spanning tree is better than Kruskal's when the graph has a large number of vertices.

True False

Question No: 26
If a problem is in NP, it must also be in $\overline{\mathrm{P}}$.
True False (p178) unknown

## Question No: 27

If a problem is in NP-complete, it must also be in NP.
True

## False

## Question No: 28

Maximum number of vertices in a Directed Graph may be $\left|V^{2}\right|$

> True False

## Question No: 29

The Huffman algorithm finds a (n) $\qquad$ solution.
Optimal (Reference) Non-optimal Exponential Polynomial

## Question No: 30

The Huffman algorithm finds an exponential solution True False (Reference)

Question No: 31
The Huffman algorithm finds a polynomial solution True False (Reference)

The codeword assigned to characters by the Huffman algorithm have the property that no codeword is the postfix of any other. True False

## Question No: 33

The codeword assigned to characters by the Huffman algorithm have the property that no codeword is the prefix of any other. True (p101) False

## Prefix Property:

The codewords assigned to characters by the Huffman algorithm have the property that no codeword is a prefix of any other:

## Question No: 34

Shortest path problems can be solved efficiently by modeling the road map as a graph. True False

## Question No: 35

Dijkstra's algorithm is operates by maintaining a subset of vertices. True False

## Question No: 36

Merge sort is stable sort, but not an in-place algorithm True (p54) False

## Question No: 37

In counting sort, once we know the ranks, we simply $\qquad$ numbers to their final positions in an output array. Delete copy (p57) Mark arrange

## Question No: 38

Dynamic programming algorithms need to store the results of intermediate sub-problems. True (p75) False

## Question No: 39

A $p \times q$ matrix A can be multiplied with a $q \times r$ matrix $B$. The result will be a $p \times r$ matrix $C$. There are ( $p . r$ ) total entries in C and each takes $\qquad$ to compute.
O (q)
O (1) (p84)
$\mathrm{O}\left(\mathrm{n}^{2}\right)$
$\mathrm{O}\left(\mathrm{n}^{3}\right)$

## Question No: 40

6. Due to left-complete nature of binary tree, heaps can be stored in Link list
Structure
Array
None of above

## Question No: 41

Which of the following is calculated with big $\mathbf{O}$ notation?
Lower bounds Upper bounds
Both upper and lower bound Medium bounds

- The definition of $\Theta$-notation relies on proving both a lower and upper asymptotic bound.
- The O-notation is used to state only the asymptotic upper bounds.


## Question No: 42

Merge sort makes two recursive calls. Which statement is true after these recursive calls finish, but before the merge step?

The array elements form a heap
Elements in each half of the array are sorted amongst themselves
Elements in the first half of the array are less than or equal to elements in the second half of the array None of the above

Question No: 43
Who invented Quick sort procedure?
Hoare Sedgewick Mellroy Coreman

## Question No: 44

Consider the following Huffman Tree. The binary code for the string TEA is:
1000010
01100010
1000110
1110110


## Question No: 45

Can an adjacency matrix for a directed graph ever not be square in shape?
Yes
No

- No. since we want to describe the relationship between each node and each other node, we need precisely $\mathrm{n}^{\wedge} 2$ matrix entries.

Question No: 46
One of the clever aspects of heaps is that they can be stored in arrays without using any $\qquad$
Pointers (p40) constants variables functions

## Question No: 47

Merge sort requires extra array storage. True (p54) False

- Mergesort is a stable algorithm but not an in-place algorithm. It requires extra array storage.


## Question No: 48

Non-optimal or greedy algorithm for money change takes $\qquad$
$\mathrm{O}(\mathrm{k}) \quad(\mathrm{p} 99) \quad \mathrm{O}(\mathrm{kN}) \quad \mathrm{O}(2 \mathrm{k}) \quad \mathrm{O}(\mathrm{N})$
Question No: 49
The Huffman codes provide a method of encoding data inefficiently when coded using ASCII standard.
True False (p99)

- The Huffman codes provide a method of encoding data efficiently.


## Question No: 50

Using ASCII standard the string "abacdaacacwe" will be encoded with $\qquad$ bits Select correct option:
96 (p101 12×8=96)

120
Question No: 51
Using ASCII standard the string abacdaacac will be encoded with $\qquad$ bits.
80
(p99) 160
320
100

- Consider the string " abacdaacac". if the string is coded with ASCII codes, the message length would be10 $\times 8=80$ bits.


## Question No: 52

Using ASCII standard the string abacdaacac will be encoded with 160 bits. True
False (p99)
Question No: 53
Using ASCII standard the string abacdaacac will be encoded with 320 bits
True
False (p99)

## Question No: 54

Using ASCII standard the string abacdaacac will be encoded with 100 bits.
True
False (p99)
Question No: 55
Using ASCII standard the string abacdaacac will be encoded with 32 bytes. True False (p99)

## Question No: 56

The greedy part of the Huffman encoding algorithm is to first find two nodes with smallest frequency.
True (p100) False

## Question No: 57

The greedy part of the Huffman encoding algorithm is to first find two nodes with character frequency. True False (p100)

## Question No: 58

The greedy part of the Huffman encoding algorithm is to first find two nodes with larger frequency.
True False
Question No: 59
Dijkestra's single source shortest path algorithm works if all edges weights are non-negative and there are negative cost cycles. True False

Question No: 60
Dijkestra s single source shortest path algorithm works if all edges weights are non-negative and there are no negative cost cycles. True (p159) False

Question No: 61
Dijkestra s single source shortest path algorithm works if all edges weights are negative and there are no negative cost cycles. True False

## Question No: 62

Floyd-Warshall algorithm is a dynamic programming algorithm; the genius of the algorithm is in the clever recursive formulation of the shortest path problem. True (p162) False

## Question No: 63

Floyd-Warshall algorithm, as in the case with DP algorithms, we avoid recursive evaluation by generating a table for $\mathrm{d}_{\mathrm{ij}}{ }^{(\mathbf{k})}$
True (p164)
False

Question No: 64
The term coloring came from the original application which was in map drawing.
True (p173) False
The term "coloring" came from the original application which was in architectural design. True
Question No: 65
In the clique cover problem, for two vertices to be in the same group, they must be $\qquad$ each other.
Apart from Far from Near to Adjacent to (p176)

## Question No: 66

In the clique cover problem, for two vertices to be in the same group, they must be adjacent to each other.
True False
Question No: 67
In the clique cover problem, for two vertices to be in the same group, they must be apart from each other.
True False (p176)
Question No: 68
The difference between Prims algorithm and Dijkstra s algorithm is that Dijkstra s algorithm uses a different key. True (p156) False

## Question No: 69

The difference between Prim s algorithm and Dijkstra s algorithm is that Dijkstra s algorithm uses a same key. True False (p156)

## Question No: 70

You have an adjective list for G , what is the time complexity to computer graph transpose $\mathrm{G}^{\wedge} \mathrm{T}$.?
(V + E ) (p138)

## Question No: 71

Given an adjacency list for $G$, it is possible to compute $G^{T}$ in $\Theta(V+E)$ time.
It takes $\mathbf{O}(\log V)$ to extract a vertex from the priority queue.
Question No: 72
Overall time for Kruskal is $\Theta(E \log E)=\Theta(E \log V)$ if the graph is sparse. (p149) True
Question No: 73
An optimization problem is one in which you want to find,
Not a solution
An algorithm
Good solution
The best solution

## Question No: 74

Although it requires more complicated data structures, Prim's algorithm for a minimum spanning tree is better than Kruskal's when the graph has a large number of vertices.

True
False

## Question No: 75

If a graph has v vertices and e edges then to obtain a spanning tree we have to delete
v edges.
$\mathrm{v}-\mathrm{e}+5$ edges
$\mathrm{v}+\mathrm{e}$ edges.
None of these

## Question No: 76

Huffman algorithm uses a greedy approach to generate an antefix code $T$ that minimizes the expected length $B$ (T) of the encoded string. True False (p102)

## Question No: 77

Huffman algorithm uses a greedy approach to generate a postfix code $T$ that minimizes the expected length $B$ (T) of the encoded string. True False (p102)

## Question No: 78

Huffman algorithm uses a greedy approach to generate a prefix code $T$ that minimizes the expected length $B(T)$ of the encoded string. True (p102) False

## Question No: 79

Bellman-Ford allows negative weights edges and negative cost cycles. True False (p159)

- Bellman-Ford allows negative weights edges and no negative cost cycles.

Question No: 80
We do sorting to,
keep elements in random positions
keep the algorithm run in linear order
keep the algorithm run in $(\log n)$ order
keep elements in increasing or decreasing order

## Question No: 81

After partitioning array in Quick sort, pivot is placed in a position such that
Values smaller than pivot are on left and larger than pivot are on right
Values larger than pivot are on left and smaller than pivot are on right
Pivot is the first element of array
Pivot is the last element of array
Question No: 82
Dynamic programming algorithms need to store the results of intermediate sub-problems. True

## Question No: 83

Which statement is true?
If a dynamic-programming problem satisfies the optimal-substructure property, then a locally optimal solution is globally optimal.

If a greedy choice property satisfies the optimal-substructure property, then a locally optimal solution is globally optimal.
Both of above None of above

## Question No: 84

What general property of the list indicates that the graph has an isolated vertex?
There is Null pointer at the end of list.
The Isolated vertex is not handled in list.
Only one value is entered in the list.
There is at least one null list.

## Question No: 85

Depth first search is shortest path algorithm that works on un-weighted graphs. True
False (p153)

- The breadth-first-search algorithm is a shortest-path algorithm that works on un-weighted graphs.


## Question No: 86

Which is true statement?
Select correct option:
Breadth first search is shortest path algorithm that works on un-weighted graphs (p153)
Depth first search is shortest path algorithm that works on un-weighted graphs.
Both of above are true.
None of above are true.
(1) In Prim's algorithm, the additional information maintained by the algorithm is the length of the shortest edge from vertex v to points already in the tree.
a) True
b) False
c) unknown
(2) Although it requires more complicated data structures, Prim's algorithm for a minimum spanning tree is better than Kruskal's when the graph has a large number of vertices.
a) True
b) False
c) unknown
(3) If a problem is NP-complete, it must also be in NP.
a) True
b) False
c) unknown
(4) Which statement is true?
(i) The running time of Bellman-Ford algorithm is T (VE)
(ii) Both Dijkstra's algorithm and Bellman-Ford are based on performing repeated relaxations
(iii) The 0-1 knapsack problem is hard to solve

Only i Only iii Both i and iii All of these
5) Which of the following arrays represent descending (max) heaps?
i. $[10,7,7,2,4,6]$
ii. [10,7,6,2,4,7]
iii. $[10,6,7,2,4,6]$

Only ii
iv. $[6,6,7,2,4,10]$

Only iv Both ii and iv Both i and iii
6. Which of the following statement(s) is/are correct?
(a) $O(n \log n+n 2)=O(n 2)$.
(b) $\mathrm{O}(\mathrm{n} \log \mathrm{n}+\mathrm{n} 2)=\mathrm{O}(\mathrm{n} 2 \log 2 \mathrm{n})$
(c) $\mathrm{O}(\mathrm{c} \mathrm{n} 2)=\mathrm{O}(\mathrm{n} 2)$ where c is a constant.
(d) $\mathrm{O}(\mathrm{c} \mathrm{n} 2)=\mathrm{O}(\mathrm{c})$ where c is a constant.
(e) $\mathrm{O}(\mathrm{c})=\mathrm{O}(1)$ where c is a constant.

Only (a) \& (e) Both (c) and (e)
7. Which of the shortest path algorithms would be most appropriate for finding paths in the graph with negative edge weights and cycles?
i.Dijkstra's Algorithm
ii. Bellman-Ford Algorithm
iii. Floyd Warshall Algorithm

Only ii $\quad$ Only iii $\quad$ Both ii \& iii
9. Suppose we have two problems A and B .Problem A is polynomial-time reducible and problem B is NPcomplete. If we reduce problem A into B then problem A becomes NP-complete Yes No
12. Edge $(\mathrm{u}, \mathrm{v})$ is a forward edge if
$u$ is a proper descendant of $v$ in the tree
$v$ is a proper descendant of $u$ in the tree
None of these
14. If, in a DFS forest of digraph $G=(V, E), f[u]=f[v]$ for an edge $(u, v)$ ? $E$ then the edge is called

Back edge Forward edge Cross Edge Tree Edge None of these
16. Best and worst case times of an algorithm may be same. True False
17. Can an adjacency matrix for a directed graph ever not be square in shape? Yes No

1. In which order we can sort?

Increasing order only decreasing order only
Increasing order or decreasing order both at the same time
3. In the analysis of Selection algorithm, we make a number of passes, in fact it could be as many as,
$\mathrm{T}(\mathrm{n}) \quad \mathrm{T}(\mathrm{n} / 2) \quad \log \mathrm{n} \quad(\mathrm{p} 37) \quad \mathrm{n} / 2+\mathrm{n} / 4$
4. How much time merge sort takes for an array of numbers?
$T\left(n^{\wedge} 2\right) \quad T(n) \quad T(\log n) \quad T(n \log n)(p 40)$
7. Sieve Technique applies to problems where we are interested in finding a single item from a larger set of $\qquad$ n items (p34) phases pointers constant
8. The sieve technique works in $\qquad$ as follows
Phases (p34) numbers integers routines
9. For the heap sort, access to nodes involves simple $\qquad$ operations.
Arithmetic binary algebraic logarithmic
10. The analysis of Selection algorithm shows the total running time is indeed $\qquad$ in n , Arithmetic geometric linear (p39) orthogonal
12. Slow sorting algorithms run in,
$\mathbf{T}\left(\mathbf{n}^{\wedge} \mathbf{2 )}(\mathbf{p 3 9 )} \quad \mathrm{T}(\mathrm{n}) \quad \mathrm{T}(\log \mathrm{n}) \quad \mathrm{T}(\mathrm{n} \log \mathrm{n})\right.$
13. A heap is a left-complete binary tree that conforms to the Increasing order only decreasing order only heap order ( $\log \mathrm{n})$ order
14. For the heap sort we store the tree nodes in Level-order traversal (p40) in-order traversal pre-order traversal post-order traversal 20: In Sieve Technique we do not know which item is of interest True (p34) False

22: Divide-and-conquer as breaking the problem into a small number of
Pivot Sieve smaller sub problems (p34) Selection
23: For the sieve technique we solve the problem,
Recursively (p34) mathematically precisely accurately
25: The reason for introducing Sieve Technique algorithm is that it illustrates a very important special case of, Divide-and-conquer (p34) decrease and conquer greedy nature 2-dimension Maxima
32. Sorting is one of the few problems where provable $\qquad$ bonds exits on how fast we can sort,
Upper Lower (p39) Average $\log n$
34: Sieve Technique can be applied to selection problem? True False
Heaps can be stored in arrays without using any pointers; this is due to the __ nature of the binary tree,
Left-complete right-complete tree nodes tree leaves
38: How many elements do we eliminate in each time for the Analysis of Selection algorithm?
n/2 elements (p36) ( $\mathrm{n} / 2$ ) +n elements
$\mathrm{n} / 4$ elements
2 n elements

42: The sieve technique is a special case, where the number of sub problems is just
5 Many
1 (p34)
few

## Question \# 1

For the Sieve Technique we take time
T(nk) (p34) $\quad T(n / 3) \quad n \wedge 2 \quad n / 3$

## Question \# 2

The number of nodes in a complete binary tree of height $h$ is
Select correct option:

$$
\begin{aligned}
& 2^{\wedge(h+1)}-1 \\
& 2 *(\mathrm{~h}+1)-1 \\
& 2 *(\mathrm{~h}+1) \\
& \left((\mathrm{h}+1)^{\wedge} 2\right)-1
\end{aligned}
$$

## Question No: 1

Random access machine or RAM is a/an
Machine build by Al-Khwarizmi
Mechanical machine
Electronics machine
Mathematical model (p10)
Question No: 2
Analysis of Selection algorithm ends up with,

T(n)
$\mathrm{T}(1 / 1+\mathrm{n})$
$\mathrm{T}(\mathrm{n} / 2)$
$\mathrm{T}((\mathrm{n} / 2)+\mathrm{n})$

## Question \# 6

Continuation sort is suitable to sort the elements in range 1 to k

1. $K$ is Large $2 . \mathrm{K}$ is not known $\quad 3 . \mathrm{K}$ may be small or large $4 . \mathrm{K}$ is small (p57)

## Question \# 5

Counting sort is suitable to sort the elements in range 1 to k :
$\begin{array}{llll}\mathrm{K} \text { is large } & \mathrm{K} \text { is small } & \mathrm{K} \text { may be large or small } & \text { None }\end{array}$

## Question \# 10

Continuing sort has time complexity of?

1. O (n)
(p58)
2. $\mathrm{O}(\mathrm{n}+\mathrm{k})$
3. $O(n \log n)$

Question \# 11
Counting sort has time complexity:
O (n) (p58) O(n+k)

## Question \# 6

Memoization is :
To store previous results for further use.
To avoid unnecessary repetitions by writing down the results of recursive calls and looking them again if needed later. (p74)
To make the process accurate
None of the above

## Question \# 9

In Quick sort algorithm, constants hidden in $T(n \lg n)$ are:
Large Medium Not known

## Question \# 12

Quick sort is based on divide and conquer paradigm; we divide the problem on base of pivot element and:
There is explicit combine process as well to conquer the solution.
No work is needed to combine the sub-arrays, the array is already sorted
Merging the subarrays
None of above.

## Question \# 13

In RAM model instructions are executed
One after another (p10)
Parallel
Concurrent
Random

## Question \# 14

In the analysis of Selection algorithm, we eliminate a constant fraction of the array with each phase; we get the convergent $\qquad$ series in the analysis,
linear arithmetic geometric (p37) exponent

Question No: 2
$\qquad$ is a graphical representation of an algorithm
$\Sigma$ notation $\quad$ Enotation Flowchart (Ref) Asymptotic notation

## Question No: 3

A RAM is an idealized machine with $\qquad$ random-access memory. 256MB $\quad 512 \mathrm{MB} \quad$ An infinitely large (p10) 100GB

## Question No: 4

What type of instructions Random Access Machine (RAM) can execute? Choose best answer
Algebraic and logic
Geometric and arithmetic
Arithmetic and logic
Parallel and recursive

## Question No: 5

What will be the total number of max comparisons if we run brute-force maxima algorithm with n elements?

$$
\begin{array}{lllll}
\mathrm{n}^{2} & (\mathrm{p} 14) & 2 \mathrm{n} / \mathrm{n} & \mathrm{n} & 8 \mathrm{n}
\end{array}
$$

## Question No: 6

What is the solution to the recurrence $T(n)=T(n / 2)+n$.
$\mathrm{O}(\operatorname{logn}) \quad \mathrm{O}(\mathrm{n}) \quad \mathrm{O}(\mathrm{n} \log n)$
$\mathrm{O}(\mathrm{n} 2)$

## Question No: 7

Consider the following code:
For $(\mathrm{j}=1 ; \mathrm{j}<\mathrm{n} ; \mathrm{j}++$ )

```
For(k=1; k<15;k++)
For(l=5;1<n;1++)
{
Do_something_constant();
\}
```

What is the order of execution for this code.
$\mathbf{O}(\mathrm{n})$
O(n3)
$\mathrm{O}(\mathrm{n} 2 \log \mathrm{n})$
$\mathrm{O}(\mathrm{n} 2)$
Question No: 8
Consider the following Algorithm:
Factorial (n) \{
if $(\mathrm{n}=1)$
return 1
else
return ( n * Factorial(n-1))
\{
Recurrence for the following algorithm is:
$\mathrm{T}(\mathrm{n})=\mathrm{T}(\mathrm{n}-1)+1$
$\mathrm{T}(\mathrm{n})=\mathrm{nT}(\mathrm{n}-1)+1$
$\mathrm{T}(\mathrm{n})=\mathrm{T}(\mathrm{n}-1)+\mathrm{n}$
$T(n)=T(n(n-1))+1$

## Question No: 9

What is the total time to heapify?
$O(\log n)(p 43) \quad O(n \log n) \quad O(n 2 \log n) \quad O(\log 2 n)$
Question No: 10
When we call heapify then at each level the comparison performed takes time
It will take $\Theta$ (1) (p43)
Time will vary according to the nature of input data
It cannot be predicted
It will take $\Theta(\log \mathrm{n})$

## Question No:

In Quick sort, we don't have the control over the sizes of recursive calls
-True (p49) -False -Less information to decide -Either true or false
Question No: 12
Is it possible to sort without making comparisons?
Yes (p57)
No

Question No: 13
If there are $\boldsymbol{\Theta}(\mathbf{n} 2)$ entries in edit distance matrix then the total running time is
$\Theta$ (1)
$\Theta(\mathrm{n} 2)$
$\Theta$ (n)
$\Theta(\mathrm{n} \log \mathrm{n})$

## Question No: 14

For Chain Matrix Multiplication we cannot use divide and conquer approach because,
We do not know the optimum $k$ (p86)
We use divide and conquer for sorting only
We can easily perform it in linear time
Size of data is not given
Question No: 15
The Knapsack problem belongs to the domain of $\qquad$ problems.
Optimization (p91) NP Complete Linear Solution Sorting

## Question No: 16

Suppose we have three items as shown in the following table, and suppose the capacity of the knapsack is 50 i.e. $\mathrm{W}=50$.
Item Value Weight
$1 \quad 60 \quad 10$
$2 \quad 100 \quad 20$
$3 \quad 120 \quad 30$
The optimal solution is to pick
Items 1 and 2
Items 1 and 3
Items 2 and 3

None of these

## Quiz 5

## Question \# 1

Theta asymptotic notation for $\mathrm{T}(\mathrm{n})$ :
Select correct option:
Set of functions described by: $\mathrm{c} 1 \mathrm{~g}(\mathrm{n})<=\mathrm{f}(\mathrm{n})$ for c 1 some constant and $\mathrm{n}=\mathrm{n} 0$
Set of functions described by $\mathrm{c} 1 \mathrm{~g}(\mathrm{n})>=\mathrm{f}(\mathrm{n})$ for c 1 some constant and $\mathrm{n}=\mathrm{n} 0$
Theta for $T(n)$ is actually upper and worst case complexity of the code
Set of functions described by: $\mathbf{c l g}(\mathrm{n})<=\mathrm{f}(\mathrm{n})<=\mathrm{c} 2 \mathrm{~g}(\mathrm{n})$ for c 1 and c 2 some constants and $\mathrm{n}=\mathrm{n} 0$

## Question \# 2

A (an) $\qquad$ is a left-complete binary tree that conforms to the heap order
Heap (page 40)
Binary tree
Binary search tree
Array

## Question \# 3

Consider the following Algorithm: Fun(n)\{ if ( $\mathrm{n}=1$ ) return 1 else return ( $\mathrm{n} * \operatorname{Fun}(\mathrm{n}-1)$ ) \} Recurrence for the above algorithm is: Select correct option:

$$
\begin{aligned}
& \mathrm{nT}(\mathrm{n}-1)+1 \\
& 2 \mathrm{~T}(\mathrm{n}-1)+1 \\
& \mathrm{~T}(\mathrm{n}-1)+\mathrm{cn} \\
& \mathrm{~T}(\mathrm{n}-1)+1
\end{aligned}
$$

## Question \# 4

The recurrence relation of Tower of Hanoi is given below $T(n)=\{1$ if $n=1$ and $2 T(n-1)$ if $n>1$ In order to move a tower of 5 rings from one peg to another, how many ring moves are required?
$\begin{array}{llllll}\text { Select correct option: } & 16 & 10 & 32 & 31\end{array}$

## Question \# 5

In Quick Sort Constants hidden in T(n $\log \mathrm{n})$ are

1. Large
2. Medium
3. Small
(Reference)
4. Not Known

## Question \# 6

In stable sorting algorithm:
One array is used
In which duplicating elements are not handled.
More than one arrays are required.
Duplicating elements remain in same relative position after sorting. (p54)

## Question \# 7

In in-place sorting algorithm is one that uses arrays for storage:
An additional array
No additional array (p54)
Both of above may be true according to algorithm
More than 3 arrays of one dimension

## Question \# 8

Which may be a stable sort?
$\begin{array}{lll}\text { 1. Merger } & \text { 2. Insertion } & \text { 3. Both above (p54) 4. None of the above }\end{array}$

## Question \# 9

An in place sorting algorithm is one that uses $\qquad$ arrays for storage

1. Two dimensional arrays
2. More than one array
3. No Additional Array
4. None of the above

## Question \# 10

Quick sort is
Stable and In place
Not stable but in place
Stable and not in place
Sometime in place and sometime stable

## Question \# 11

One Example of in place but not stable sort is
Quick (p54) Heap Merge Bubble

## Question \# 12

Which may be stable sort?
Bubble sort Insertion sort both of above (p54)

## Question \# 13

Merge sort is stable sort, but not an in-place algorithm True (p54) False
An in-place sorting algorithm is one that uses no additional array for storage. A sorting algorithm is stable if duplicate elements remain in the same relative position after sorting.
$9|3| 3^{\prime}|5| 6\left|5^{\prime}\right| 2|1| 3^{\prime \prime} \quad$ unsorted

| $1\|2\| 3\left\|3^{\prime}\right\| 3^{\prime \prime}\|5\| 5^{\prime}\|6\| 9$ | stable sort |
| :--- | :--- |


| $1\|2\| 3^{\prime}\|3\| 3^{\prime \prime}\left\|5^{\prime}\right\| 5\|6\| 9$ | unstable |
| :--- | :--- |

Bubble sort, insertion sort and selection sort are in-place sorting algorithms.
Bubble sort and insertion sort can be implemented as stable algorithms but selection sort cannot (without significant modifications).

Mergesort is a stable algorithm but not an in-place algorithm. It requires extra array storage.
Quicksort is not stable but is an in-place algorithm.
Heapsort is an in-place algorithm but is not stable.

